Processes and Threads

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Policy vs. Mechanism

Policy – Decisions about what should be done.

 Mechanism – Algorithms and data structures that *implement* policy

Processes

- Process \rightarrow program in execution
- Processes
 - are assigned resources
 - run in user or kernel mode
- Things we do with resources
 - share
 - acquire
 - release

Multiprogramming

- Multiple processes on a single computer
- Scheduling algorithm selects which processes are allocated CPU
- P₃₂ P₂₅
 Don't know when scheduled next... avoid loops for timing



Multiprogramming

- How do we transition from one process to another?
 - cooperative multitasking processes explicitly yield



 preemptive multitasking forced release of processor resource based on external conditions



Process life cycle

- initialization created by another process
- running / blocked / ready – Process is active

- termination
 - Voluntary
 - error
 - complete



- Involuntary
 - unhandled exception
 - killed by another process

Process State









Queuing diagram

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Process Implementation

- Process control block (PCB)
 - kernel data structure representing processes
 - frequently implemented as a fixed size array
- Process control block contains
 - state
 - what else?

Typical PCB

Process managementMeRegistersPoiProgram counterPoiProgram status wordPoiStack pointerProcess statePriorityScheduling parametersProcess IDParent processProcess groupSignalsTime when process startedCPU time usedCPU time usedChildren's CPU timeTime of next alarmImagement	nory management ater to text segment info ater to data segment info ater to stack segment info	File management Root directory Working directory File descriptors User ID Group ID
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Process types

Background – no user supervision

Interactive – User I/O



 Daemon – Special background processes that provide services



Program thread

- Thread of execution
 - Stream of instructions being executed
 - CPU registers
 - Stack
 - current procedure calls
 - local variables

Can we have more than one thread?

Multithreading

- Multiple threads (aka lightweight processes)
- Threads within process share resources:
 - files, heap, and any other allocated resources
 - and are allocated the CPU, much like processes
- Thread control block
 - Keeps registers, PC, PSW and state

Multithreading dangers

- Threads can
 - overwrite each others' stacks
 - access data structures in transient states
 - change heap values
 - access resources in unexpected interleaving



Lost in Space, Irwin Allen Productions

Okay, so why bother?

Threads are convenient



Threads are convenient

A multithreaded Web server.



Tanenbaum 2008

Creating a thread

Implementation dependent

- POSIX implementation (man pages/FAQ for details)
 - headers: <pthread.h> <sched.h>
 - pthread_create(...) Create a thread
 - sched_yield(...) Next thread runs
 - pthread_exit(...) Terminate thread
 - pthread_join(...) Wait for specific thread to exit
 - pthread_attr_init(...) Initialize options structure to be passed to pthread_create

POSIX thread example

```
int main(int argc, char *argv[])
{
     /* The main program creates 10 threads and then exits. */
     pthread_t threads[NUMBER_OF_THREADS];
     int status, i;
     for(i=0; i < NUMBER_OF_THREADS; i++) {
          printf("Main here. Creating thread %d0, i);
          status = pthread_create(&threads[i], NULL, print_hello_world, (void *)i);
          if (status != 0) {
                printf("Oops. pthread_create returned error code %d0, status);
                exit(-1);
     exit(NULL);
}
```

See the course FAQ for a concrete example.

POSIX thread example

 Previous example needs header: #include <pthread.h>

 Compilation on a linux box gcc –o threadeg threadeg.c –I pthread –I rt

Note: gcc puts library flags *after* list of files (rarely done) and the order of libraries is important.

Thread mechanisms

- user-level threads
 - implemented via a user library
 - scheduling occurs in user code
- kernel-level threads
 - part of OS implementation
 - data structures are maintained in kernel code

User- vs. Kernel- mode threads



User- vs. Kernel- mode threads

- Kernel-mode
 - Kernel schedules threads, not processes
 - Thread operations and switches require shift to kernel mode (\$\$\$)
- User-mode
 - Kernel unaware of user threads
 What happens when one thread blocks?
 - Very fast thread operations

Hybrid thread design

- User-threads mapped onto kernel-threads
- Best of both worlds



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Pop-up threads

Dynamic creation of threads to handle events



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Threads

 Suppose a thread calls a function that sets a global return code (e.g. UNIX errno)

• Can we run into problems?